

Quantum computation: lecture 10 (Ruediger)

Shor's algorithm: summary

Task: given N , non-prime, find a non-trivial factor of N

We assume :

- 2 does not divide N

- $N \neq p^e$, p prime

(as these cases are easily solvable)

Algorithm: Pick $a \in \{2..N-1\}$ unif. at random

- Compute $\gcd(a, N) = d$ with Euclid's algo
- if $d \neq 1$, we have found a non-trivial factor of N
- if $d = 1$, compute the multiplicative order of $a^2 \pmod N$, call it r (NB: r is the smallest value s.t. $a^r \pmod N = 1$) \rightarrow Shor
- if r is odd, then declare failure
- if r is even, write $a^r - 1 = (a^{r/2} - 1) \cdot (a^{r/2} + 1)$

- if N divides $a^{\frac{F}{2}} + 1$, then declare failure
- otherwise, compute $\gcd(N, a^{\frac{F}{2}} - 1)$ and $\gcd(N, a^{\frac{F}{2}} + 1)$: both must be non-trivial: done

Analysis $\rightarrow P(\text{failure}) \leq \frac{1}{4}$ (repeat the algo)
(in this case)

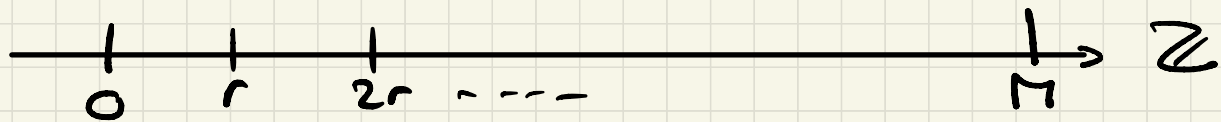
Take now $M = 2^m \sim N^2$ ($r < N$).

We wish to compute the period of $f: \mathbb{Z} \rightarrow \mathbb{Z}$

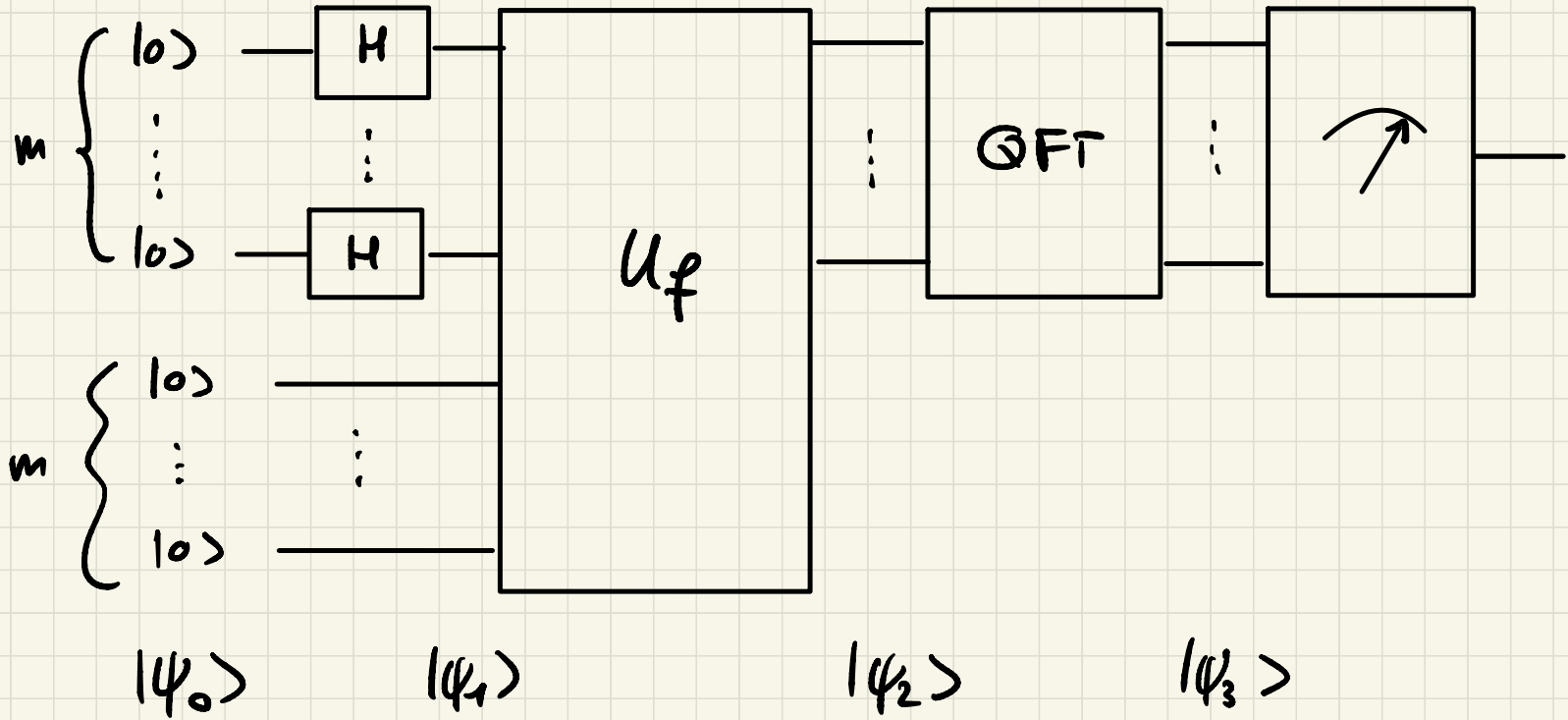
defined as $f_{a,N}(x) = a^x \pmod N$

f is r -periodic: $f(x+r) = f(x) \quad \forall x \in \mathbb{Z}$

Look at f on $\{0 \dots M-1\}$



Consider now the following quantum circuit:



$$|\psi_0\rangle = \underbrace{|0\dots 0\rangle}_m \otimes \underbrace{|0\dots 0\rangle}_m$$

$$|\psi_1\rangle = \frac{1}{\sqrt{M}} \sum_{x=0}^{M-1} |x\rangle \otimes |0\dots 0\rangle$$

$$|\psi_2\rangle = \frac{1}{\sqrt{M}} \sum_{x=0}^{M-1} |x\rangle \otimes |f(x)\rangle$$

[Assume r divides M]

$$= \frac{1}{\sqrt{M}} \sum_{x_0=0}^{r-1} \sum_{j=0}^{\frac{M}{r}-1} |x_0 + jr\rangle \otimes \underbrace{|f(x_0 + jr)\rangle}_{= f(x_0)}$$

$$|\psi_3\rangle = \frac{1}{M} \sum_{x_0=0}^{r-1} \sum_{y=0}^{M-1} e^{\frac{2\pi i x_0 y}{M}} \sum_{j=0}^{\frac{M}{r}-1} e^{\frac{2\pi i j y}{M/r}} |y\rangle \otimes |f(x_0)\rangle$$

Measurement: $P_y = |y\rangle\langle y| \otimes I_m$

Outcome:

If r divides M , then $P(y) = \frac{1}{M^2} \sum_{x_0=0}^{r-1} \left| \sum_{j=0}^{\frac{M}{r}-1} e^{\frac{2\pi i j y}{M/r}} \right|^2$

So $y = k \cdot \frac{M}{r}$ with $k \in \{0 \dots r-1\}$ uniformly dist.

Now, if r does not divide M , the formula becomes:

$$P(y) = \frac{1}{M^2} \sum_{x_0=0}^{r-1} \left| \sum_{j=0}^{A(x_0)-1} e^{\frac{2\pi i j y}{M/r}} \right|^2 \quad \text{where } A(x_0) = \begin{cases} \text{either } \lceil \frac{M}{r} \rceil \\ \text{or } \lceil \frac{M}{r} \rceil - 1 \end{cases}$$

In this second case, it can be shown that

$$P\left(\exists 0 \leq k \leq r-1 \text{ st. } \left|y - k \cdot \frac{M}{r}\right| \leq \frac{1}{2}\right) \geq \frac{2}{5}$$

i.e. $\left|\frac{y}{M} - \frac{k}{r}\right| \leq \frac{1}{2M} \leq \frac{1}{2r^2}$ in case of success.

Assume now $\gcd(k, r) = 1$ (this happens with probability $\geq \frac{1}{4 \ln \ln r}$).

Systematic way to find r :

Compute all convergents of $\frac{y}{N}$ (continued fractions)

\Rightarrow look at all the denominators:

each of these denominators is a candidate

for the period r : check if $a^r \bmod N = 1$

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