

ÉCOLE POLYTECHNIQUE FÉDÉRALE DE LAUSANNE
School of Computer and Communication Sciences

Foundations of Data Science
Fall 2022

Assignment date: Friday, February 3rd, 2023, 9:15 am
Due date: Friday, February 3rd, 2023, 12:15 noon

Final Exam – SG0211

This exam is open book. No electronic devices of any kind are allowed. There are 4 problems. Choose the ones you find easiest and collect as many points as possible. We do not necessarily expect you to finish all of them. Good luck!

Name: _____

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Problem 1 (*Fisher Goes Exponential*). [15 pts]

Let $p_\theta(x)$ denote a family of distributions parameterized by θ . Define the Fisher information as

$$I_\theta = \mathbb{E}_\theta[\nabla_\theta \log p_\theta(X)(\nabla_\theta \log p_\theta(X))^T].$$

- (1) [5pts] Let $p_\theta(x) = h(x)e^{(\theta, \phi(x)) - A(\theta)}$ be an exponential family. What is the Fisher information in terms of the parameters of the family?
- (2) [5pts - 1pt per question] Consider distributions of the form $p_\lambda(x) = \lambda e^{-\lambda x}$, where $\lambda \in \mathbb{R}^+$.
 1. Write it in the form of an exponential family.
 2. What is $\Theta = \{\theta \in \mathbb{R} : A(\theta) < \infty\}$.
 3. Is the family regular?
 4. Is it minimal?
 5. What is the Fisher information?
- (3) [5pts - 1pt per question] Consider distributions of the form $p_p(k) = (1 - p)^k p$, where $p \in (0, 1)$ and $k \in \mathbb{N}$.
 1. Write it in the form of an exponential family.
 2. What is $\Theta = \{\theta \in \mathbb{R} : A(\theta) < \infty\}$.
 3. Is the family regular?
 4. Is it minimal?
 5. What is the Fisher information?

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Problem 2 (*Compression*). [15 pts]

Suppose $\mathcal{P} \in \Pi(\mathcal{X}, \mathcal{Y})$ be a probability distribution on $\mathcal{X} \times \mathcal{Y}$ and (X, Y) be a joint random variable with distribution P_{XY} with marginals P_X and P_Y .

In what follows, assume that **all codes are optimal, prefix-free, and binary**. Optimal here means having smallest possible average length. All logs are to the base 2.

- (1) [1 pt] Let $c_X : \mathcal{X} \rightarrow \{0, 1\}^*$ and $c_Y : \mathcal{Y} \rightarrow \{0, 1\}^*$ be optimal prefix free codes. What are lower and upper bounds for the expected length of these codes c_X and c_Y ?
- (2) [1 pt] Let $c_{XY} : \mathcal{X} \times \mathcal{Y} \rightarrow \{0, 1\}^*$ be an optimal prefix free code. What are lower and upper bounds for the expected length of this code?
- (3) [10 pts total] In this sub problem, assume that X, Y have a joint distribution according to the following table:

	Y=0	Y=1
X=0	1/4	0
X=1	1/8	1/8
X=2	1/8	1/8
X=3	0	1/4

- (a) [4 pts] What are lower and upper bounds for the expected lengths of c_X and c_Y ? Are the lower bounds tight?
 - (b) [3 pts] What are lower and upper bounds for the expected lengths of c_{XY} ? Is the lower bound tight?
 - (c) [3 pts] For the above joint distribution, is it more efficient to compress separately and concatenate the individual code words (which, as we saw in the lecture, is guaranteed to yield a prefix free code), or to compress (X, Y) jointly (again, in a prefix free manner)?
- (4) [3 pts] Assume that (X, Y) has some generic joint distribution. Assume further that $I(X; Y) > 1$. Show that in this case optimal joint prefix free compression is more efficient than compressing individually and concatenating.

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Problem 3 (*Stability implies Generalization*). [12 pts]

Let $S = \{(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n)\}$ be a training dataset composed of n i.i.d. samples drawn from \mathcal{D} . As usual, we denote $L_{\mathcal{D}}(h) = E_{(x,y) \sim \mathcal{D}}[l(h(x), y)]$ and $L_S(h) = \frac{1}{n} \sum_{i=1}^n l(h(x_i), y_i)$ the true and empirical risks of a hypothesis h , respectively. For simplicity, let us denote by h_S the output of a learning algorithm when trained with dataset S .

An important property of learning algorithms is their ability to generalize, i.e., the true and empirical risks of the output hypothesis should be close in expectation. Formally, we say that a learning algorithm \mathcal{A} ϵ -generalizes in expectation if

$$|E_S[L_S(h_S) - L_{\mathcal{D}}(h_S)]| < \epsilon. \quad (1)$$

An interesting connection arises when we investigate the *stability* of a learning algorithm. Formally, we call a learning algorithm ϵ -uniformly stable if $\forall S, S'$ datasets of size n that differ in at most one sample we have

$$\sup_{(x,y)} l(h_S(x), y) - l(h_{S'}(x), y) < \epsilon. \quad (2)$$

Notations: $(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n), (\tilde{x}_1, \tilde{y}_1), \dots, (\tilde{x}_n, \tilde{y}_n)$ are $2n$ independently sampled training examples. We define $S = \{(x_1, y_1), \dots, (x_n, y_n)\}$, $\tilde{S} = \{(\tilde{x}_1, \tilde{y}_1), \dots, (\tilde{x}_n, \tilde{y}_n)\}$ and $S^{(i)} = \{(x_1, y_1), \dots, (x_{i-1}, y_{i-1}), (\tilde{x}_i, \tilde{y}_i), (x_{i+1}, y_{i+1}), \dots, (x_n, y_n)\}$.

(1) [2 pts] Prove that $L_{\mathcal{D}}(h_S) = E_{\tilde{S}}[\frac{1}{n} \sum_{i=1}^n l(h_S(\tilde{x}_i), \tilde{y}_i)]$.

(2) [3 pts] Prove that $E_{S, \tilde{S}}[l(h_S(\tilde{x}_i), \tilde{y}_i)] = E_{S, S^{(i)}}[l(h_{S^{(i)}}(x_i), y_i)]$.

- (3) [7 pts] Prove that an ϵ -uniformly stable learning algorithm ϵ -generalizes in expectation, by justifying each step in the following sequence.

$$\begin{aligned}
|E_S[L_S(h_S) - L_{\mathcal{D}}(h_S)]| &\stackrel{(a)}{=} \left| E_S \left[L_S(h_S) - E_{\tilde{S}} \left[\frac{1}{n} \sum_{i=1}^n l(h_S(\tilde{x}_i), \tilde{y}_i) \right] \right] \right| \\
&\stackrel{(b)}{=} \left| E_S [L_S(h_S)] - E_{S, \tilde{S}} \left[\frac{1}{n} \sum_{i=1}^n l(h_S(\tilde{x}_i), \tilde{y}_i) \right] \right| \\
&\stackrel{(c)}{=} \left| E_S [L_S(h_S)] - \frac{1}{n} \sum_{i=1}^n E_{S, \tilde{S}} [l(h_S(\tilde{x}_i), \tilde{y}_i)] \right| \\
&\stackrel{(d)}{=} \left| E_S [L_S(h_S)] - \frac{1}{n} \sum_{i=1}^n E_{S^{(i)}, (x_i, y_i)} [l(h_{S^{(i)}}(x_i), y_i)] \right| \\
&\stackrel{(e)}{=} \left| E_S \left[\frac{1}{n} \sum_{i=1}^n l(h_S(x_i), y_i) \right] - \frac{1}{n} \sum_{i=1}^n E_{S, S^{(i)}} [l(h_{S^{(i)}}(x_i), y_i)] \right| \\
&\stackrel{(f)}{=} \left| \frac{1}{n} \sum_{i=1}^n E_{S, S^{(i)}} [l(h_S(x_i), y_i) - l(h_{S^{(i)}}(x_i), y_i)] \right| \\
&\stackrel{(g)}{\leq} \frac{1}{n} \sum_{i=1}^n \epsilon = \epsilon
\end{aligned}$$

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Problem 4 (*Multi-arm Bandits*). [20 pts]

We consider the following game where in each round t we can choose between $[N] = \{1, 2, \dots, N\}$ different actions. After we choose an action $a_t \in [N]$ an adversary reveals the loss of each action in this round, call it $l_i^t \in [0, 1]$, $i \in [N]$. Note that this is an adversarial setting, where the losses do not come from a probability distribution. This setting differs from what we had discussed in class where only the loss for the chosen action was revealed.

Our goal is to design a randomized algorithm \mathcal{A} which maintains a probability distribution p^t over actions, and achieves a sub-linear regret, i.e., $\mathcal{R}(T) = \max_i \{ \sum_{t=1}^T E_{A_t \sim p^t} [l_{A_t}^t - l_i^t] \} \leq o(T)$. We also note that the adversary may know the probability distribution p^t , but does not know the realizations A_t . We will analyze the following algorithm:

Algorithm 1: Multiplicative Weights Update

Input: learning parameter ϵ

Initialization: $p_i^1 = 1/N, w_i^1 = 1, \forall i \in [N], \Phi^1 = N$

for $t = 1$ to T **do**

$A_t \sim p^t$

 Adversary reveals the loss vector l^t and we suffer $l_{A_t}^t$

 Update weights $w_i^{t+1} = w_i^t \cdot \exp(-\epsilon \cdot l_i^t), \forall i \in [N]$ and let $\Phi^{t+1} = \sum_i w_i^{t+1}$

 Update the probability distribution: $p_i^{t+1} = w_i^{t+1} / \Phi^{t+1}, \forall i$

end for

- (1) [2 pts] Prove that $w_i^{T+1} = \exp(-\epsilon \cdot \sum_{t=1}^T l_i^t), \forall i \in [N]$
- (2) [8 pts] Prove that $\Phi^{t+1} \leq \Phi^t \cdot \exp(\epsilon^2 - \epsilon \langle p^t, l^t \rangle)$
Hint: Note that $w_i^{t+1} = p_i^{t+1} \cdot \Phi^{t+1}$ and use the inequalities: (a) $e^x \leq 1 + x + x^2, \forall x \in [0, 1]$ and (b) $e^x \geq x + 1, \forall x$.
- (3) [2 pts] Prove that $\Phi^{T+1} \leq \Phi^1 \cdot \exp(\epsilon^2 \cdot T - \epsilon \sum_{t=1}^T \langle p^t, l^t \rangle)$
- (4) [8 pts] By noting that $\Phi^1 \cdot \exp(\epsilon^2 \cdot T - \epsilon \sum_{t=1}^T \langle p^t, l^t \rangle) \geq \Phi^{T+1} \geq w_i^{T+1}, \forall i \in [N]$ set the learning parameter ϵ so that $\mathcal{R}(T) \leq 2\sqrt{\log(N) \cdot T}$.

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